Errata to ASSAT: Computing Answer Sets of A Logic Program By SAT Solvers *Artificial Intelligence* 157(1-2): 115-137 (2004).

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Definition 1 about loops on page 119 in the paper was not correct. The following is a corrected version:

> Given a finite normal logic program P that may contain constraints, let $atom(P)$ be the set of atoms in it, and G_P its positive dependency graph. A non-empty subset L of $atom(P)$ is called a loop of P if for any p and q in L, there is a path of length > 0 from p to q in G_P such that all the vertices in the path are in L.

This means that if L is non-empty, and not a singleton, then L is a loop if and only if the subgraph of G_P induced by L is strongly connected. If L is a singleton, say $\{p\}$, then it is a loop if and only if there is an arc from p to p in G_P. If L is a singleton but not a loop, then one could still construct its "loop formula", which will in fact be a consequence of Clark's completion of the program.[1] Thus Theorem 1, the main theorem of the paper, still holds if we generalize loops to include every singleton set of atoms.

[1]We are grateful to Joohyung Lee and Vladimir Lifschitz for discussions that made us aware of this.