

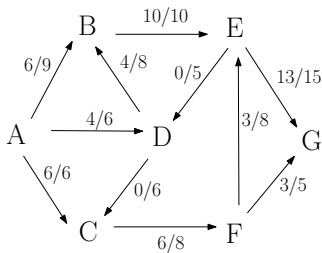
COMP3711H: Design and Analysis of Algorithms

Tutorial 10
Revision of November 21, 2016

HKUST

Question 1

Consider the given graph with flow values f and capacities c (f/c) as shown. $s = A$ and $t = G$.



Draw, the residual graph.

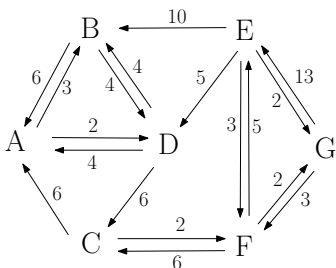
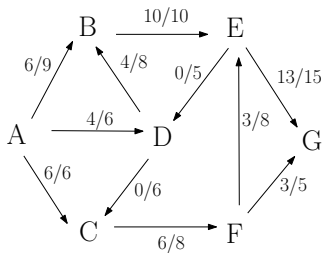
Find an augmenting path.

Show the new flow created by adding the augmenting path flow. Is your new flow optimal?

Prove or disprove.

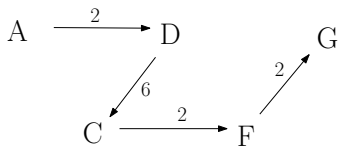
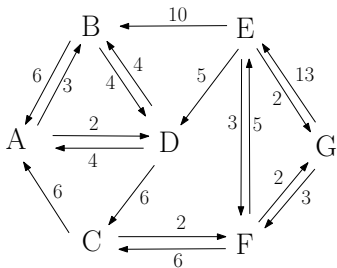
Solution 1

Draw, the residual graph.



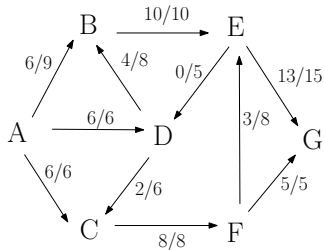
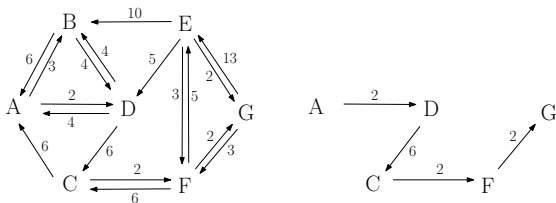
Solution 1

Find an augmenting path.



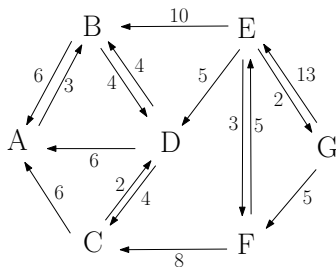
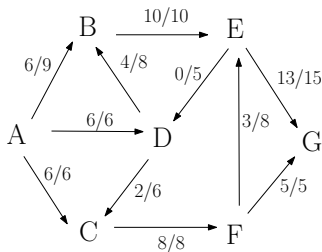
Solution 1

Show the new flow created by adding the augmenting path flow.



Solution 1

Is your new flow optimal? Prove or disprove.

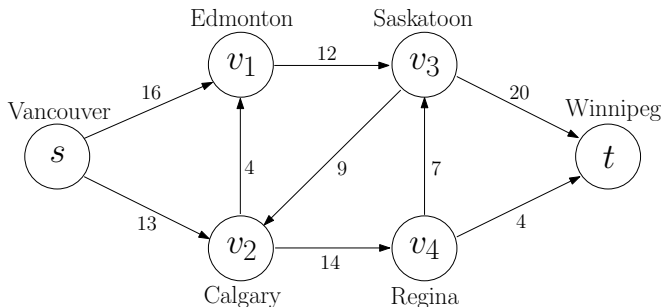


Yes it is. Consider the S, T cut defined by $S = \{A, B, C, D\}$ which is the set of nodes that can be reached from A in the residual graph of the new flow. Then $T = \{E, F, G\}$. Note that the capacity of the cut is $C(S, T) = 10 + 8 = 18$ which is exactly the value of the flow leaving A .

\Rightarrow from the max-flow min-cut theorem, this is a maximal flow.

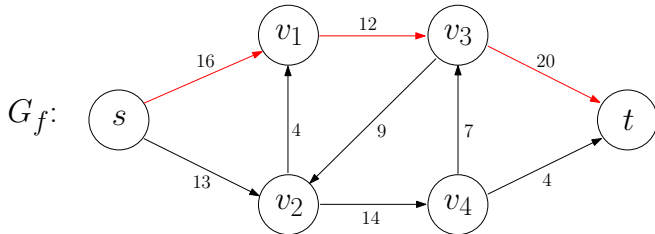
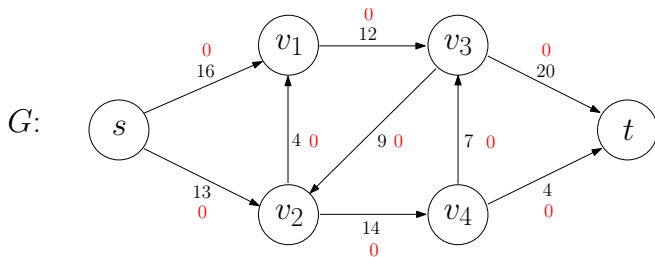
Question 2

Show the execution of the Edmonds-Karp algorithm on the following flow network.

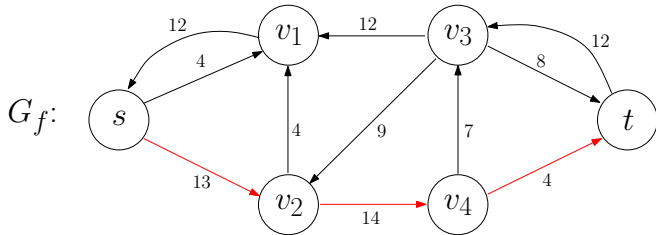
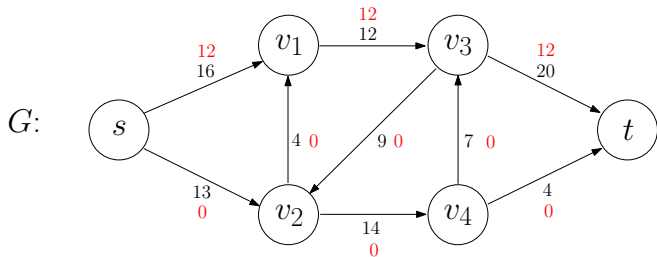


Recall that the *Edmonds-Karp* Algorithm is the version of the Ford-Fulkerson method that always chooses a shortest (by number of edges) Augmenting Path.

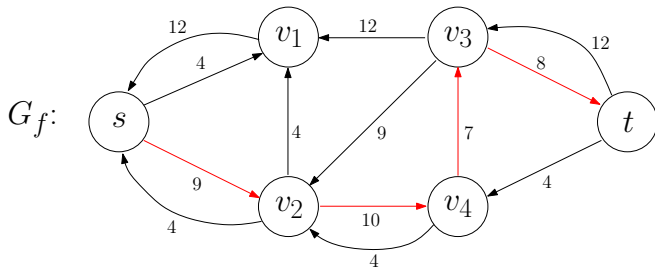
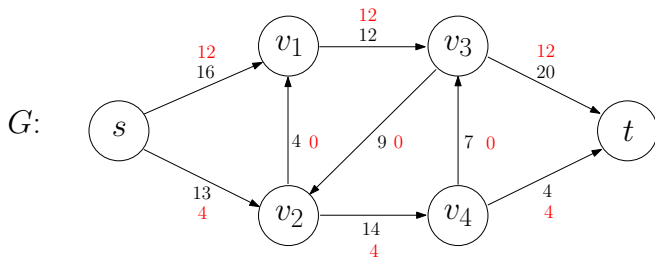
Solution 2



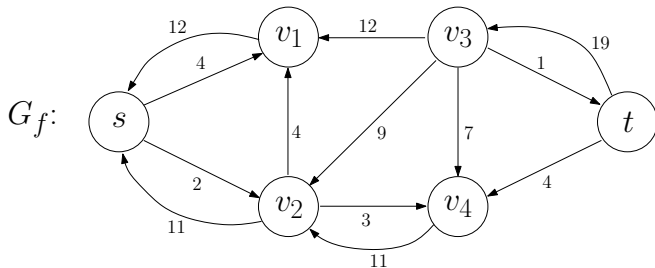
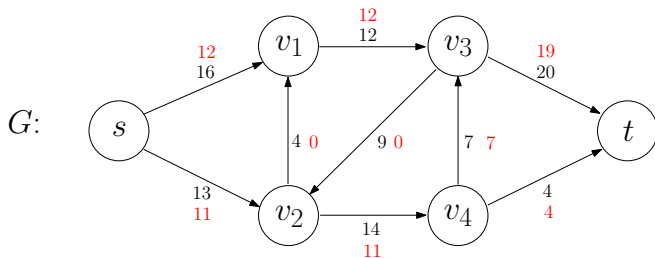
Solution 2



Solution 2



Solution 2



Solution 2

